

# Pleading for a functorial approach of Delzant correspondence

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# Preliminary comments

- This is a *work in progress* with Damien Lejay : it is possible that despite our research in the literature, our results are already known, or that the question get a disappointing answer. One goal for this talk is to check that with the audience.
- Category theory / functorial approach is essentially a language: "*Its merit is that it exists*". It can help identify the core of a given problem, not make it disappear !
- Classification of almost-toric systems "*à la Delzant*" is the definition and the study of its moduli space. The only definition of a moduli space comes from algebraic geometry, and is written in the language of category. Defining what is "the" category of almost-toric systems, starting with "the" category of toric systems shall help us understand the meaning of "*à la Delzant*" beyond some heuristic approach.

# Plan

Joint work with D. Lejay.

- 1 Reminder of the language of categories
  - Definition
  - Functors and equivalence of categories
- 2 The existing categorie(s) of toric systems and Delzant polytopes
  - On hamiltonian torus action
  - Toric integrable systems
  - The Delzant correspondence
  - A category of toric systems
- 3 Can we do better ?
  - Polarized toric varieties

## Definition

A (small) category  $\mathcal{C}$  is defined as

- a set of objects  $Ob(\mathcal{C})$ ,
- For each pair of objects  $A$  and  $B$  of  $\mathcal{C}$ , a set  $\mathcal{C}(A, B)$  called the set of “morphisms”, or arrows :  $f : A \rightarrow B$ .
- For each triplet  $A, B, C$  of objects in  $\mathcal{C}$ , a binary operation

$$\circ : \mathcal{C}(A, B) \times \mathcal{C}(B, C) \rightarrow \mathcal{C}(A, C)$$

$$(f, g) \mapsto \circ(f, g) =: g \circ f$$

that is associative :  $(f \circ g) \circ h = f \circ (g \circ h)$ , and with an identity morphism  $id_A \in \mathcal{C}(A, A)$  for each object in  $\mathcal{C}$  :  $f \circ id_A = f$ ,  $id_A \circ g = g$ .

## Example

- **Set** :  $Ob(\mathbf{Set}) = \{\text{all sets}\}$ , and  $\mathbf{Set}(X, Y) := \{\text{maps from } X \text{ to } Y\}$
- **Top** :  $Ob(\mathbf{Top}) = \{\text{topological spaces}\}$ , and  
 $\mathbf{Top}((X, \tau), (Y, \tau')) := \{\text{continuous maps from } (X, \tau) \text{ to } (Y, \tau')\}$
- **Grp** :  $Ob(\mathbf{Grp}) = \{\text{groups}\}$  and  
 $\mathbf{Grp}(G, H) := \{\text{group homomorphisms from } G \text{ to } H\}$
- A groupoid is a category for which every arrow has an inverse.

## Example (**Rel**)

- $Ob(\mathbf{Rel}) = \{\text{all sets}\}$
- $\mathbf{Rel}(A, B) :=$   
 $\{\text{all binary relations between } A \text{ and } B \text{ i.e. : subsets of } A \times B\}$ ,
- for  $\mathcal{R} : A \rightarrow B$  and  $\mathcal{R}' : B \rightarrow C$ ,  $\mathcal{R}' \circ \mathcal{R}$  is defined by :  $z(\mathcal{R}' \circ \mathcal{R})x$  if there exists a  $y \in B$  such that  $y\mathcal{R}x$  and  $z\mathcal{R}'y$ .

## Definition

A functor  $F : \mathcal{C} \rightarrow \mathcal{C}'$  associates objects of  $\mathcal{C}$  to objects of  $\mathcal{C}'$ ,  $A \rightarrow F(A)$ , and to each morphism  $f \in \mathcal{C}(A, B)$ , associates a morphism  $F(f) \in \mathcal{C}'(F(A), F(B))$  such that :

$$F(f \circ_{\mathcal{C}} g) = F(f) \circ_{\mathcal{C}'} F(g) \text{ and } F(id_A) = id_{F(A)}$$

## Example

- $\pi_1 : \mathbf{Top}_p \rightarrow \mathbf{Grp}$  which associates to a pointed topological set  $(X, p)$  its fundamental group  $\pi_1(X, p)$
- $Spec : \mathbf{ComRing} \rightarrow \mathbf{LocRngSp}$ , which associates to a commutative ring its spectrum of prime ideals. An important property in algebraic geometry is that it defines an **equivalence of category** between  $\mathbf{ComRing}^{op}$  and the “image” of  $Spec$ , which are called *affine schemes*.

# Some “mantras” in category theory

- “*Morphisms are everything*”: more than the objects, they give the “shape” of the category  $\Rightarrow$  Adjusting the number of morphisms is like having enough open sets in a topology
- “*Do NOT make choices*”: the power of category language is to formulate construction as “best/universal solution to a problem” which is formulated by a diagram.  
Each non-canonical choice hinders the power, and thus the purpose of using category, so we try to avoid them as much as possible.

## Definition

A torus is here a compact connected abelian Lie group of finite dimension, here denoted  $\mathbf{T}$ .  $\ker(\exp)$  defines a canonical lattice  $\Lambda$  on  $\mathfrak{t}$ , and by duality  $\Lambda^*$  on  $\mathfrak{t}^*$ .

Let  $(M^{2n}, \omega)$  be a symplectic manifold.

## Definition

An Lie group action  $\rho : G \rightarrow \text{Diff}(M)$  is called Hamiltonian if there exists a Lie algebra homomorphism  $\eta : \mathfrak{g} \rightarrow \mathcal{C}^\infty(M, \mathbb{R})$  called the *comoment map* such that

$$\begin{array}{ccc}
 \mathfrak{g} & \xrightarrow{d\rho} & \Gamma(M, TM)^G \\
 & \searrow \eta & \uparrow ad \\
 & & \mathcal{C}^\infty(M)
 \end{array}$$

The moment map  $\mu$  is defined by duality. Both are uniquely defined up to a constant.



# First definition of an Integrable Hamiltonian systems

## Definition

An **Integrable Hamiltonian System** (IHS) will be a triplet  $(M^{2n}, \omega, F)$  with  $(M, \omega)$  a symplectic manifold, and  $F := (f_1, \dots, f_n) \in \mathcal{C}^\infty(M^{2n} \rightarrow \mathbb{R}^n)$ , such that:

- $\forall i, j = 1..n, \{f_i, f_j\}_\omega = 0,$
- $\text{rk}(dF)$  is maximal almost everywhere.

$F$  is the moment map for a Poisson  $\mathbb{R}^n$ -action,

$\mathcal{F} := \{\text{c.c. of } F^{-1}(c) \mid c \in F(M)\}$  its associated (singular) Lagrangian foliation with projection  $\pi_{\mathcal{F}} : M \rightarrow B$  so that  $\mathcal{F} = \{\Lambda_b\}_{b \in B}$ .

For us,  $M$  will always be compact connected.

We set  $B_r := \{b \in B \mid \forall m \in \Lambda_b, \text{rk}(dF_m) = n\}$

# Action-Angles coordinates

## Theorem (Liouville-Arnold-Mineur theorem)

Let  $b \in B_r$  the set of regular value of  $\pi_{\mathcal{F}}$ , then there exists an open set  $\mathcal{U}$  of  $\Lambda_b$ , and an  $n$ -dimensional torus  $\mathbf{T}$  such that  $\mathcal{F}|_{\mathcal{U}}$  is symplectically isomorphic to a fibration by Lagrangian tori in a saturated neighborhood of the zero section in  $T^*\mathbf{T}$ .

$\pi_{\mathcal{F}}$  defines a *singular* Lagrangian torus fibration.

**A byproduct :** At each  $b \in B_r$ , the integral covectors define *canonical* lattice  $A_b$  on  $T_b^*B$ :  $(B_r, \mathcal{A} := (A_b)_{b \in B_r})$  is an open integral affine manifold:  $GL_n(\mathbb{Z}) \times \mathbb{R}^n$ .

$B$  and  $F(M)$  are related by:

On  $B_r$  (at least) the map  $\tilde{F} : (B, \mathbf{C}) \rightarrow (F(M), \mathcal{C}^\infty(\cdot, \mathbb{R}))$  is a local diffeomorphism.

$$\begin{array}{ccc}
 (M^{2n}, \omega) & \xrightarrow{\pi_{\mathcal{F}}} & B \supseteq B_r \\
 & \searrow F & \downarrow \tilde{F} \\
 & & F(M)
 \end{array}$$

$\simeq$

# To chose or not to chose ?

## Definition

An *intrinsic*, or *geometric* integrable system is a singular Lagrangian torus fibration  $p : M \rightarrow B$ .

An intrinsic integrable system is called “genuine” if there exists an  $F$  such that  $p = \pi_{\mathcal{F}}$ ; an *immersed* integrable system is an intrinsic integrable system together with a choice of an  $F$ .

## Definition

A genuine system  $(M, \omega, \pi_{\mathcal{F}})$  is called toric if there exists an (effective) action of a torus whose moment map is a choice of  $F$ .

An “*immersed*” toric system is a genuine toric system together with a choice of such an action.

The connectedness of the fiber in [AGS] ensures that in the toric case, the immersion is actually an embedding.

## Theorem (Atiyah – Guillemin & Sternberg, 1982)

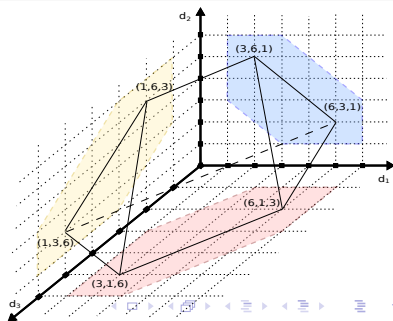
Let  $\mathbf{T} \curvearrowright (M^{2n}, \omega)$  be an Hamiltonian action of a torus of dimension  $d$ , with moment map  $\mu : M^{2n} \rightarrow \mathfrak{t}^*$ . We have :

- the fibers of  $\mu$  are connected,
- $\mu(M) =: \Delta$  is an intrinsic polytope i.e. : an integral affine **convex** manifold with corners with a finite number of extremal points, that is genuine, ie. there exists a global section for  $\mathcal{A}_\Delta$ .

### Example

The Gelfand-Cetlin system with  $\Lambda = (1, 3, 6)$ : the set of isospectral  $3 \times 3$ -Hermitian matrices have diagonals  $(d_1, d_2, d_3)$  contained in the polytope

$$\Delta := \text{Conv}((1, 3, 6), (1, 6, 3), \dots)$$



## Theorem (Delzant '87)

- Let  $(M, \omega, \pi_{\mathcal{F}})$  be an intrinsic toric integrable system. Then: *Delta* is moreover normal and smooth
- Given two “embedded” toric integrable system manifolds, **with the same torus**  $T$ . If  $\mu_1(M_1) = \mu_2(M_2)$ , there exists a  $T$ -equivariant symplectomorphism such that the following diagram commutes

$$\begin{array}{ccc}
 M_1 & \xrightarrow{\Phi} & M_2 \\
 \mu_1 \downarrow & & \downarrow \mu_2 \\
 \mu_1(M_1) & \equiv & \mu_2(M_2)
 \end{array}$$

Such an symplectomorphism is uniquely defined up to automorphisms.

- Given a Delzant polytope  $\Delta$ , there is an explicit construction for a canonical genuine toric integrable system. If  $\Delta$  is given together with a choice of an embedding, one can reconstruct an “embedded” toric integrable system, which is unique up to a  $T$ -equivariant symplectomorphism.

From that, one can give definition for the category of genuine and embedded toric systems, and the category of genuine and embedded Delzant polytopes

## Definition

We set

- $\mathbf{gSysTor}_0$  the trivial category whose object are genuine toric systems, and with identity morphisms only,
- $\mathbf{eSysTor}_0$  the category whose object are embedded toric systems, and the morphisms are the  $T$ -equivariant symplectomorphisms,
- $\mathbf{gDel}_0$  the trivial category whose object are genuine Delzant polytopes, and with the identity morphisms only,
- $\mathbf{eDel}_0$  the trivial category whose object are embedded Delzant polytopes, and with the identity morphisms only,

Not satisfying from a categorical viewpoint: **very poor categories !**

$\mathbf{eSysTor}_0$  is a groupoid; the Delzant classification defines an equivalence of categories, but it is just a bijection between their  $\pi_0$ 's.

We need more maps, that are not isomorphisms !

**A hint:** Arnol'd-Liouville -Mineur theorem

$\implies$  the integral affine structure is the crucial structure !  $\mathbf{gDel}_0$  is a (non-full !)  
!) subcategory of the category of integral affine manifolds.

$\implies$  A possible  $\mathbf{gDel}_1$  with

$gDel_1(\Delta, \Delta') = \{a : \Delta \rightarrow \Delta' \mid a(\Lambda) \text{ a sublattice of } \Lambda'\} ?$

l. e. :  $\mathbf{eDel}_1$  with  $eDel_1(\Delta, \Delta') = \{A \in GL_n(\mathbb{Z}) \times \mathbb{R}^n \mid A(\Delta) \subseteq \Delta'\} ?$

# The underlying toric variety

There exists a forgetful functor to the category of toric varieties.

## Theorem

*There are no non-constant algebraic map from  $\mathbb{C}P^n$  to any algebraic variety of smaller dimension.*

Proof : Bézout theorem !



# Morphisms of fans

## Definition

A morphism  $D : \sigma \rightarrow \sigma'$  is a integral linear map such that every cone  $\tau \subseteq \sigma$  is sent to a cone  $\tau' \subseteq \sigma'$

## Definition

A morphism of toric varieties  $V \rightarrow W$  is the same as a morphism of fan from  $\sigma_V \rightarrow \sigma_W$

$$\begin{array}{ccc}
 M & \xrightarrow{\Phi} & \\
 \mu_1 \downarrow & & \downarrow \mu_2 \\
 \mu_1(M_1) & \longleftarrow \! \! \longleftarrow & \mu_2(M_2)
 \end{array}$$

Let  $(M^{2m}, \omega_M, \mathbf{T}, \rho_M, \mu_M)$  and  $(N^{2n}, \omega_N, \mathbf{S}, \rho_N, \mu_N)$  two embedded toric systems.

## Definition

Let  $\varphi \in \mathcal{C}^\infty(M, N)$  and  $\psi$  be such that  $\varphi^* \omega_N = \omega_M$ , and such that it is  $(S, T)$ -equivariant, i.e : the following diagram commutes

$$\begin{array}{ccc}
 T \times M & \xrightarrow{id_T \times \varphi} & S \times N \\
 \rho_M^T \downarrow & & \downarrow \rho_N^S \circ (\psi \times id_M) \\
 M & \xrightarrow{\varphi} & N
 \end{array}$$



# Thank you !