Numerical Aspects of Special Functions

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- Can we rely on Maple or Mathematica ?
- Writing a book on the numerical aspects of special functions

- For the first time after the 1964 version a complete revision of Abramowitz and Stegun
- The new DLMF (Digital Library of Mathematical Functions) will appear in a hardcover edition and as a free electronic publication on the World Wide Web.
- New chapters on mathematical methods (such as computer algebra, asymptotic analysis) and new functions classes (such as *q*-hypergeometric functions, functions of matrix argument) will be included.
- This project started at NIST in 1999. It is expected now that the new book, the web version, and the CD version become available in 2006.



What to have at my desert island?







Figure 2: The Airy function |Ai(z)| in the complex plane



Figure 3: The Hankel function $|H_0^{(1)}(z)|$ in the complex plane



To compute

- Bessel functions
- Legendre functions
- Confluent hypergeometric functions (Kummer, Whittaker, Coulomb)
- Parabolic cylinder functions
- Gauss hypergeometric functions
- Incomplete gamma and beta functions
- Orthogonal polynomials

recursion relations are important and frequently used.



The recursion relations are the form

$$y_{n+1} + a_n y_n + b_n y_{n-1} = 0, \quad n = 1, 2, 3, \dots$$

When there are two linearly independent solutions f_n and g_n , such that

$$\lim_{n \to \infty} \frac{f_n}{g_n} = 0,$$

the computation of the minimal solution

$$f_n, \quad n = 2, 3, 4, \dots$$

from f_0 and f_1 (forward recursion) is usually very unstable. The solution g_n is called a dominant solution.

For example, the functions

$$f_n(x) = e^x - 1 - x - \dots - \frac{x^{n-1}}{(n-1)!}, \quad f_0(x) = e^x, \quad f_1(x) = e^x - 1,$$

satisfy the recursion relation

$$y_{n+1} - \left(1 + \frac{x}{n}\right)y_n + \frac{x}{n}y_{n-1} = 0, \quad n = 1, 2, 3, \dots$$

Computing $f_n(1)$ with Maple, standard 10 Digits, we see that

$$f_{13}(x) = -0.340710500 \times 10^{-9},$$

a negative number.



The function

$$g_n(x) = e^x - f_n(x) = 1 + x + \dots + \frac{x^{n-1}}{(n-1)!},$$

is for x > 0 a dominant solution of the recursion. It can be computed in a stable way with starting values

$$g_0(x) = 0, \quad g_1(x) = 1.$$

The recursion for the functions $f_n(x)$ follows from the simpler recursion

$$y_{n+1} = y_n - \frac{x^n}{n!}.$$

Use this in backward direction with false starting value $f_{21}(1) = 0$. Then

$$f_{13}(1) = 1.7287667139 \times 10^{-10} \dots,$$

which is correct in all shown digits.

Backward recursion for a minimal solution with false starting values is the basis for the Miller algorithm.



The Gauss hypergeometric function is defined by

$${}_2F_1\left(\begin{array}{c}a,\ b\\c\end{array};z\right) = \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n n!} z^n, \quad (a)_n = \frac{\Gamma(a+n)}{\Gamma(a)}.$$

In this form it is defined for |z| < 1, and $c \neq 0, -1, -2, \ldots$.

Several recursion relations with respect to the parameters a, b, c exist. In fact, we have 26 relations for the functions

$$y_n = {}_2F_1 \left(\begin{matrix} a + \varepsilon_1 n, \ b + \varepsilon_2 n \\ c + \varepsilon_3 n \end{matrix}; z \right),$$

where $\varepsilon_j = -1, 0, 1$ (not all $\varepsilon_j = 0$).



For example, the Gauss hypergeometric function

$$y_n = {}_2F_1\left(\begin{array}{c}a+n, \ b\\c\end{array}; z\right)$$

satisfies the recursion relation

$$(a+n)(z-1)y_{n+1} + (2a+2n-c-az-nz+bz)y_n + (c-a-n)y_{n-1} = 0.$$

There are important questions (from a numerical point of view).

Questions:

- Is this relation stable for computing this Gauss function? That is, can we compute $y_n, n \ge 2$, from y_0, y_1 in a stable way, or is y_n a minimal solution?
- Can we find a second solution of this relation, and, if yes, is this solution a minimal or dominant solution?
- Have these solutions the same properties (minimal, dominant) for all complex values of z?

We consider recursion with respect to n of the hypergeometric function

$${}_{2}F_{1}\left(\begin{matrix}a+\varepsilon_{1}n, b+\varepsilon_{2}n\\c+\varepsilon_{3}n\end{matrix}; z\right)$$

where

$$\varepsilon_j = -1, 0, \text{ or } 1.$$

There are $3^3 - 1 = 26$ possible recursions.



Because of the symmetry relation

$${}_{2}F_{1}\left(\begin{array}{c}a,\ b\\c\end{array};\ z\right) = {}_{2}F_{1}\left(\begin{array}{c}b,\ a\\c\end{array};\ z\right)$$

and the functional relation

$$_{2}F_{1}\begin{pmatrix}a, b\\c\ \end{pmatrix} = (1-z)^{-a}{}_{2}F_{1}\begin{pmatrix}a, c-b\\c\ \end{bmatrix} \frac{z}{z-1}$$

only 5 basic forms need to be studied.

The remaining 21 recursions follow from these 5 recursions.



The basic forms are

$${}_{2}F_{1}\begin{pmatrix}a+n, b+n\\c & ;z\end{pmatrix}, {}_{2}F_{1}\begin{pmatrix}a+n, b+n\\c-n & ;z\end{pmatrix},$$
$${}_{2}F_{1}\begin{pmatrix}a+n, b\\c & ;z\end{pmatrix}, {}_{2}F_{1}\begin{pmatrix}a+n, b\\c-n & ;z\end{pmatrix},$$
$${}_{2}F_{1}\begin{pmatrix}a, b\\c+n & ;z\end{pmatrix}.$$

It is also important to consider the following three cases

$$_{2}F_{1}\left(a, b \\ c-n; z \right),$$

$${}_{2}F_{1}\left(\begin{array}{c}a-n,\ b\\c+n\end{array};z\right),$$
$${}_{2}F_{1}\left(\begin{array}{c}a-n,\ b-n\\c+n\end{array};z\right),$$

which follow from the five basic forms by changing the n-direction.

Remember the theory of the Gauss differential equation. The six functions

$$w_1 = {}_2F_1\begin{pmatrix}a, b\\c\\c\end{pmatrix}, \qquad w_2 = z^{1-c}{}_2F_1\begin{pmatrix}1+a-c, 1+b-c\\2-c\\c\end{pmatrix},$$

$$w_{3} = {}_{2}F_{1} \begin{pmatrix} a, b \\ a+b+1-c \end{pmatrix}, \qquad w_{4} = (1-z)^{c-a-b} {}_{2}F_{1} \begin{pmatrix} c-a, c-b \\ c+1-a-b \end{pmatrix}, \qquad (a+b+1-c)^{c-a-b} {}_{2}F_{1} \begin{pmatrix} c-a, c-b \\ c+1-a-b \end{pmatrix},$$

$$w_5 = (z^{-1}e^{i\pi})^a {}_2F_1\left(\begin{array}{c}a, \ a+1-c\\a+1-b\end{array}; \frac{1}{z}\right), \quad w_6 = (z^{-1}e^{i\pi})^b {}_2F_1\left(\begin{array}{c}b, \ b+1-c\\b+1-a\end{aligned}; \frac{1}{z}\right)$$

satisfy the same differential equation.

Each w_j can be written as a linear combination of two other w-functions. For example

$$w_1 = \frac{\Gamma(a+1-c)\Gamma(b+1-c)}{\Gamma(1-c)\Gamma(a+b+1-c)}w_3 - \frac{\Gamma(a+1-c)\Gamma(b+1-c)\Gamma(c-1)}{\Gamma(1-c)\Gamma(a)\Gamma(b)}w_2.$$

Observations:

- Each term in the right-hand side satisfies the same differential equation as w_1 .
- Each term in the right-hand side satisfies the same difference equation as w_1 .
- Hence, by considering several linear combinations, we can find several solutions of the same difference equation (recursion relation).



Given the recursion relation

$$y_{n+1} + a_n y_n + b_n y_{n-1} = 0$$

we compute

$$\alpha = \lim_{n \to \infty} a_n, \quad \beta = \lim_{n \to \infty} b_n,$$

and consider the characteristic polynomial

$$t^2 + \alpha t + \beta = 0$$

with zeros t_1 and t_2 .

Then, if $|t_1| \neq |t_2|$, the difference equation has two linear independent solutions f_n and g_n with the properties

$$\frac{f_{n+1}}{f_n} \sim t_1, \quad \frac{g_{n+1}}{g_n} \sim t_2.$$

If $|t_1| > |t_2|$, we call f_n a dominant solution, and g_n the minimal solution.



For all hyper recursions, t_1 and t_2 do not depend on a, b, c; they may depend on z.

For all hyper recursions, t_1 and t_2 are non-zero and finite (except for some simple *z*-values).

By verifying when $|t_1| = |t_2|$, we find domains in the complex *z*-plane where we have to identify minimal and dominant solutions of the five basic recursion relations.

The recursion for $_2F_1\left(egin{array}{c} a+n,\ b+n\\ c \end{array};z
ight).$ There is a pair $\{f_n,g_n\}$ with

$$\lim_{n \to \infty} \frac{f_{n+1}}{f_n} = t_1 = \frac{1}{(1 - \sqrt{z})^2},$$

$$\lim_{n \to \infty} \frac{g_{n+1}}{g_n} = t_2 = \frac{1}{(1 + \sqrt{z})^2},$$

The equation $|t_1| = |t_2|$ holds when $z \le 0$. The function related with w_3 is minimal; the other five functions related with the other w_j are dominant.



In this case we have

$$y_1 = {}_2F_1 \left(\begin{array}{c} a+n, \ b+n \\ c \end{array}; z \right),$$

$$y_3 = \frac{\Gamma(a+n+1-c)\Gamma(b+n+1-c)}{\Gamma(a+b+2n+1-c)} {}_2F_1 \begin{pmatrix} a+n, \ b+n \\ a+b+2n+1-c \end{pmatrix}; \ 1-z \end{pmatrix}.$$

When z is not on the negative real axis: y_1 is dominant, y_3 is minimal.

The Jacobi polynomial is a special case:

$$P_n^{(\alpha,\beta)}(x) = \binom{n+\alpha}{n} \left(\frac{1+x}{2}\right)^{\alpha+\beta+n+1} {}_2F_1\left(\begin{array}{c} \alpha+1+n, \ \alpha+\beta+n+1\\ \alpha+1 \end{array}; \frac{x-1}{x+1} \right)$$

Forward recursion is not unstable for $P_n^{(\alpha,\beta)}(x)$ when $x \in [-1,1]$.

The recursion for
$$_2F_1\left(egin{array}{c}a+n,\ b+n\\c-n\ \end{array};z
ight).$$

There is a pair $\{f_n,g_n\}$ with

$$\lim_{n \to \infty} \frac{f_{n+1}}{f_n} = t_1 = \frac{32(1+w)}{(3+w)^3},$$

$$\lim_{n \to \infty} \frac{g_{n+1}}{g_n} = t_2 = \frac{32(1-w)}{(3-w)^3},$$

where $w = \sqrt{1 + 8z}$. The equation

$$|t_1| = |t_2|$$

holds on the curve shown in Figure 4.



Figure 4: The curve $|t_1| = |t_2|$ for the case (+ + -).

The function y_2 related with w_2 is minimal in the domain inside the curve; the function related with w_3 is minimal in the domain outside the curve; the other functions related with the other w_i are dominant.

The corresponding y_1 (always dominant), y_2 , y_3 for this case:

$$y_1 = {}_2F_1 \begin{pmatrix} a+n, b+n \\ c-n \end{pmatrix},$$

$$y_2 = (-z)^n \frac{\Gamma(b+1-c+2n)\Gamma(a+1-c+2n)}{\Gamma(a+n)\Gamma(b+n)\Gamma(1-c+n)\Gamma(2-c+n)} \times$$

$$_{2}F_{1}\left(\begin{matrix} 1+a-c+2n,\ 1+b-c+2n\\ 2-c+n \end{matrix}; z \end{matrix}\right),$$

$$y_{3} = \frac{\Gamma(b+1-c+2n)\Gamma(a+1-c+2n)}{\Gamma(1-c+n)\Gamma(a+b+1-c+3n)} {}_{2}F_{1} \begin{pmatrix} a+n, b+n \\ a+b+1-c+3n \end{pmatrix}; 1-z \end{pmatrix}.$$
The recursion for
$$_2F_1\left(egin{array}{c} a+n,\ b\\ c \end{array};z
ight)$$

There is a pair $\{f_n,g_n\}$ with

$$\lim_{n \to \infty} \frac{f_{n+1}}{f_n} = t_1 = 1,$$

$$\lim_{n \to \infty} \frac{g_{n+1}}{g_n} = t_2 = \frac{1}{1-z}$$

The equation

$$|t_1| = |t_2|$$

holds on the curve shown in Figure 5, a circle with centre 1 and radius 1.







Figure 5: The curve $|t_1| = |t_2|$ for the case (+00).

The function related with w_3 is minimal inside the circle; the function related with w_5 is minimal outside the circle; the other functions related with the other w_j are dominant.

The y_1 (always dominant), y_3 , y_5 for this case:

$$y_1 = {}_2F_1 \left(\begin{array}{c} a+n, \ b \\ c \end{array}; z \right),$$

$$y_{3} = \frac{\Gamma(a+n+1-c)}{\Gamma(a+b+n+1-c)^{2}} F_{1} \begin{pmatrix} a+n, b \\ a+b+n+1-c \end{pmatrix}; 1-z \end{pmatrix},$$

$$y_5 = (-z)^{-n} \frac{\Gamma(n+1-c+a)}{\Gamma(1+a+n-b)^2} F_1 \left(\begin{array}{c} a+n, \ a+n+1-c\\ a+n+1-b \end{array}; \frac{1}{z} \right)$$

The recursion for $_2F_1\left(egin{array}{c}a+n,\ b\\c-n\ \end{array};z
ight).$ There is a pair $\{f_n,g_n\}$ with

$$\lim_{n \to \infty} \frac{f_{n+1}}{f_n} = t_1 = 1,$$

$$\lim_{n \to \infty} \frac{g_{n+1}}{g_n} = t_2 = -\frac{4z}{(1-z)^2}.$$

The equation

$$|t_1| = |t_2|$$

holds on the curve shown in Figure 6.





Figure 6: The curve $|t_1| = |t_2|$ for the case (+0-).

The function related with w_2 is minimal in the domain inside the inner curve; the function related with w_3 is minimal in the domain between the two curves; the function related with w_5 is minimal in the domain outside the outer curve; the other functions related with the other w_j are dominant.



The y_1 (always dominant), y_2 , y_3 , y_5 for this case:

$$y_1 = {}_2F_1\left(\begin{array}{c}a+n,\ b\\c-n\end{array};z\right),$$

$$y_2 = (-z)^n \frac{\Gamma(b+1-c+n)\Gamma(a+1-c+2n)}{\Gamma(a+n)\Gamma(1-c+n)\Gamma(n+2-c)} {}_2F_1 \begin{pmatrix} 1+a-c+2n, \ 1+b-c+n \\ 2-c+n \end{pmatrix}; z \end{pmatrix},$$

$$y_3 = \frac{\Gamma(b+1-c+n)\Gamma(a+1-c+2n)}{\Gamma(1-c+n)\Gamma(a+b+1-c+2n)} {}_2F_1 \begin{pmatrix} a+n, \ b \\ a+b+1-c+2n \end{pmatrix}; 1-z \end{pmatrix},$$

$$y_5 = (-z)^{-n} \frac{\Gamma(1+a-c+2n)}{\Gamma(1-c+n)\Gamma(1+a-b+n)} {}_2F_1 \begin{pmatrix} a+n, a+1-c+2n \\ a+1-b+n \end{pmatrix}; \frac{1}{z} \end{pmatrix}.$$

The recursion for $_2F_1\begin{pmatrix}a, b\\c+n\\ \end{pmatrix}$. There is a pair $\{f_n, g_n\}$ with

$$\lim_{n \to \infty} \frac{f_{n+1}}{f_n} = t_1 = 1, \quad \lim_{n \to \infty} \frac{g_{n+1}}{g_n} = t_2 = \frac{z-1}{z}.$$

The equation $|t_1| = |t_2|$ holds when $\Re z = \frac{1}{2}$. The function related with w_1 is minimal when $\Re z < \frac{1}{2}$; the function related with w_4 is minimal when $\Re z > \frac{1}{2}$; the other functions related with the other w_j are dominant.

The y_1 , y_4 for this case:

$$y_1 = {}_2F_1\left(\begin{array}{c}a, \ b\\ c+n\end{array}; z\right),$$

$$y_4 = (z-1)^n \frac{\Gamma(n+c)}{\Gamma(c+1-a-b+n)} \times$$

$$_{2}F_{1}\left(\begin{matrix} c+n-a, \ c+n-b\\ c+n+1-a-b \end{matrix}; 1-z \end{matrix}\right)$$



The (0 0 -), (-0 +), (- - +) recursions:

These recursions need extra attention, although they are "negative n" cases of some of the basic forms.

The domains for these recursions are the same as those for the $(0\ 0\ +), (+\ 0\ -), (+\ +\ -)$ cases, respectively, which we have done earlier.

However, for each case (and each domain) the minimal and dominant solutions have to be identified again.

Conclusions (so far):

- We have identified minimal and dominant solutions for all 26 recursion relations for the hypergeometric functions.
- We have described the domains in the complex z-plane where these minimal and dominant solutions have been identified.
- Proofs of these properties (not mentioned in this lecture) are available, and are based on behaviour of solutions near the singular points $0, 1, \infty$ of the Gauss differential equation.



Literature:

- Papers on recursions (theory): Wong & Li (1992a,b).
- Papers on recursions (numerics): Gautschi (1967), Olver (1967).
- Book on recursions: Wimp (1984).
- Books on asymptotics: Olver, Wong, Luke.
- Recent papers on asymptotics of Gauss functions: Jones, Olde Daalhuis, Temme.
- Paper on numerics of Gauss functions: Forrey (1997).
- Project on recursion of Kummer functions: Deaño & Segura (2005).
- Papers on recursion of Legendre functions: Gil & Segura (1997, 1998, ...).
- Two papers on recursion of Gauss hypergeometric functions: Gil, Segura, Temme (JCAM, ??).

Consider the following questions.

- 1. Can we compute the Gauss hypergeometric function by using power series and the many transformations that are available for this function?
- 2. Can we compute this function for all possible values of *a*, *b* and *c*?

- Power series
- Integrals
- Differential equation
- Recursion relations
- Chebychev expansions
- Continued fractions

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The standard power series is

$$_{2}F_{1}\begin{pmatrix}a, b\\c\ \end{pmatrix} = \sum_{n=0}^{\infty} \frac{(a)_{n}(b)_{n}}{(c)_{n} n!} z^{n}, \quad |z| < 1.$$

Together with relations such as

$$_{2}F_{1}\begin{pmatrix}a, b\\c\ \end{pmatrix} = (1-z)^{-a}{}_{2}F_{1}\begin{pmatrix}a, c-b\\c\ \end{bmatrix}; \frac{z}{z-1}$$

For this function with need, when we use its power series,

$$\left|\frac{z}{z-1}\right| < 1.$$

In several other ways we can write the Gauss function in terms of other Gauss functions. The new arguments are

$$\frac{1}{z}$$
, $1-z$, $\frac{1}{1-z}$, $\frac{z}{z-1}$, $\frac{z-1}{z}$.

For numerical computations we need convergence conditions like

$$\begin{aligned} |z| < \rho, \ \left|\frac{1}{z}\right| < \rho, \ |1-z| < \rho, \ \left|\frac{1}{1-z}\right| < \rho, \ \left|\frac{z}{z-1}\right| < \rho, \ \left|\frac{z-1}{z}\right| < \rho, \end{aligned}$$
 with
$$0 < \rho < 1.$$



Can we cover the whole z—plane with some number ρ ? In the green domain one of the conditions

 $|z| < \rho, \left|\frac{1}{z}\right| < \rho, \left|1-z\right| < \rho, \left|\frac{1}{1-z}\right| < \rho, \left|\frac{z}{z-1}\right| < \rho, \left|\frac{z-1}{z}\right| < \rho$ is satisfied. This is the case $\rho = 0.5$.





This is the case $\rho = 0.75$.

In the yellow domains, 'around' the points $e^{\pm \frac{1}{3}\pi i}$, the standard power series cannot be used with this value of ρ .



Other pitfalls:

- For large values of a or b instabilities arise, and the power series converge slowly.
- For certain combinations of a, b and c removable singularities occur.

Example c = a + b:

$${}_{2}F_{1}\left(\begin{array}{c}a,\ b\\c\end{array};\ z\right) = \frac{\Gamma(c)\Gamma(c-a-b)}{\Gamma(c-a)\Gamma(c-b)}{}_{2}F_{1}\left(\begin{array}{c}a,\ b\\a+b-c+1\end{array};\ 1-z\right) +$$

$$\frac{\Gamma(c)\Gamma(a+b-c)}{\Gamma(a)\Gamma(b)}{}_2F_1\left(\begin{array}{c}c-a,\ c-b\\c-a-b+1\end{array};\ 1-z\right)$$

This relation is well defined for c = a + b because singularities are cancelled.

If $c \sim a + b$ numerical computations are not stable.

Summarizing so far:

- power series are very efficient in certain domains of the complex plane
- not all z-values can be covered
- instabilities occur for certain values of the parameters.

Numerical quadrature for special functions

The standard integral representations of special functions are not always suitable for numerical computations.

When parameters are large, integrals with oscillatory integrands can be very unstable representations.

By using complex contours for these integrals, or for transformed versions, stable representations can be obtained.

Numerical quadrature to compute special functions

The next task is selecting a suitable quadrature rule for computations.

Because the integrands are analytic functions, high-precision quadrature rules can be selected.

Numerical quadrature to compute special functions

Example: Bessel function ($h = \pi/n$, x = 5).

$$\pi J_0(x) = \int_0^\pi \cos(x \sin t) \, dt = h + h \sum_{j=1}^{n-1} \cos\left[x \sin(h \, j)\right] + R_n,$$

n	R_n
4	$12 \ 10^{-0}$
8	$48 \ 10^{-6}$
16	$11 \ 10^{-21}$
32	$13 \ 10^{-62}$
64	$13 \ 10^{-163}$
128	$53 \ 10^{-404}$

Much better than the standard estimates of R_n . Explanation: periodicity and smoothness.



Consider

$$F(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda\sqrt{t^2 + 1}} dt.$$

• Maple 9.5, Digits = 10, for $\lambda = 10$, gives

F(10) = -.1837516481 + .5305342893i.



Consider

$$F(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda\sqrt{t^2 + 1}} dt.$$

• Maple 9.5, Digits = 10, for $\lambda = 10$, gives

F(10) = -.1837516481 + .5305342893i.

• With Digits = 40, the answer is

F(10) = -.1837516480532069664418890663053408790017 +

+ 0.5305342892550606876095028928250448740020 i.



Consider

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+0.5305342892550606876095028928250448740020i.

So, the first answer seems to be correct in all shown digits.

Take another integral, which is almost the same:

$$F(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda\sqrt{t^2 + 1}} dt \quad \Longrightarrow \quad G(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda t} dt.$$

▶ Maple 9.5, Digits=10, for $\lambda = 10$, gives $G(10) = -0.1257674520 \times 10^{-15}$.

Take another integral, which is almost the same:

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- Maple 9.5, Digits=10, for $\lambda = 10$, gives $G(10) = -0.1257674520 \times 10^{-15}$.
- With Digits = 40, the answer is $G(10) = .16 \times 10^{-43}$.

Take another integral, which is almost the same:

$$F(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda\sqrt{t^2 + 1}} dt \implies G(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda t} dt.$$

- Maple 9.5, Digits=10, for $\lambda = 10$, gives $G(10) = -0.1257674520 \times 10^{-15}$.
- With Digits = 40, the answer is $G(10) = .16 \times 10^{-43}$.
- The correct answer is $G(\lambda) = \sqrt{\pi}e^{-\lambda^2}$ and for $\lambda = 10$ we have $G(10) = 0.6593662989 \times 10^{-43}$.

The message is: one should have some feeling about the computed result.

Otherwise a completely incorrect answer can be accepted.

Mathematica is more reliable here, and says:

"NIntegrate failed to converge to prescribed accuracy after 7 recursive bisections in t near t = 2.9384615384615387".
By the way, Maple 7 could do the following integral

$$H(\lambda) = \int_{-\infty}^{\infty} e^{-t^2 + 2i\lambda\sqrt{t^2}} dt,$$

and the funny answer was, after some simplification,

$$H(\lambda) = \sqrt{\pi}e^{-\lambda^2} [1 + \operatorname{signum}(t) \operatorname{erf} i\lambda],$$

where $\operatorname{erf} z$ is the error function. In Maple 9.5 the answer is

$$H(\lambda) = \sqrt{\pi}e^{-\lambda^2}(1 + i \text{erf } \lambda).$$

Consider

$$F(u) = \int_0^\infty e^{uit} \frac{dt}{t - 1 - i}, \quad u > 0.$$

Numerical quadrature gives F(2) = -0.934349 - 0.70922i. Mathematica 4.1 gives for u = 2 in terms of the Meijer G-function:

$$F(2) = \pi G_{2,3}^{2,1} \begin{pmatrix} 0, \frac{1}{2} \\ 0, 0, \frac{1}{2} \end{pmatrix}; 2 - 2i \end{pmatrix}.$$

Mathematica evaluates: F(2) = -0.547745 - 0.532287i.

Ask Mathematica to evaluate F(u):

$$F(u) = e^{iu - u} \Gamma(0, iu - u).$$

This gives F(2) = -0.16114 - 0.355355i.

So, we have three numerical results:

$$F_1 = -0.934349 - 0.70922i,$$

$$F_2 = -0.547745 - 0.532287i,$$

$$F_3 = -0.16114 - 0.355355i.$$

Observe that $F_2 = (F_1 + F_3)/2$. F_1 is correct.

Maple 9.5:

$$F(u) = e^{iu-u} \operatorname{Ei}(1, iu - u) = e^{iu-u} \Gamma(0, iu - u),$$

same as Mathematica. This is a wrong answer.

Next, Maple 9.5, with u = 2,

$$F(2) = e^{2i-2} \operatorname{Ei}(1, 2i-2) + 2\pi i e^{2i-2},$$

giving F(2) = -.9343493872 - .7092195102i, which is the correct answer.



A book on numerics of special functions

The topics mentioned in this lecture, and several other topics, will be discussed extensively, with examples of software, in a new book with the title of this talk.

Written together with my co-authors Amparo Gil and Javier Segura (Santander, Spain).

The project is not finished yet and the publication date is not known yet.

To be published by SIAM.